

Kronecker product optimization in SPOT

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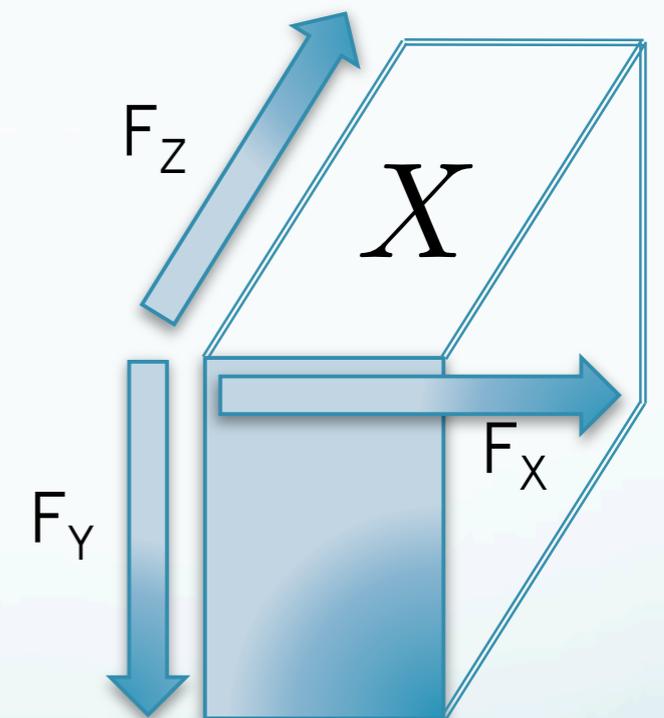
Introduction

➤ Kronecker product use

- Multidimensional transforms on sets of data
- Example: n -D Fourier transform F

$$F(X) = (F_n \otimes F_{n-1} \otimes \dots \otimes F_1) X(:,)_{\text{dim1}} \\ (+\text{final reshaping})$$

where : $X(:,)_{\text{dim1}}$ data set X collapsed along dimension 1



Introduction

➤ **Kronecker product interest**

- KP decomposition **induces weak dependences** on multiplied data (presence of block diagonal operators)
 - possibility of parallel calculations
 - optimization analyzing data's selected parts

➔ **Possible speeding-up of computations**

Program

I. Initial Kronecker product (KP)

1. General approach
2. KP multiplication algorithm in SPOT
3. Initial result

II. KP optimization

1. Permutations
2. Replacement of the *for* loop
3. Final comparison

I. Initial Kronecker product

1. General approach

➤ Definition

$$A \otimes B = \begin{pmatrix} a_{11}B & \dots & a_{1c_A}B \\ \vdots & \ddots & \vdots \\ a_{r_A1}B & \dots & a_{r_Ac_A}B \end{pmatrix}$$

➤ KP representation in SPOT

- $K = \text{opKron}(A, B)$
- ... or with as many operators

as wished: $K = \text{opKron}(A_1, A_2, \dots, A_n)$

1. General approach

- SPOT KP is an **implicit operator**:
 - never evaluated directly
 - only defined by its interaction with others operators when multiplying them
 - works on the selected parts of the right-hand multiplied operator

$$(A \otimes B)X = \text{KP}(A, B)X$$

In Matlab

$$(A \otimes B)X = \text{Function}(A, B, X)$$

In SPOT

➔ **faster than Matlab's KP multiplication on a vector**

1. General approach

Example:

`A=rand(80);`

`B= rand(80);`

`X=rand(1600,1);`

- In **MATLAB**: **1.67s** on average
- In **SPOT**: **0.03s** on average

→ **56 times faster than Matlab's KP**

2. KP multiplication algorithm in SPOT

➤ Being given X a vector: $(A \otimes B)X = (A \otimes I_{r_B})(I_{c_A} \otimes B)X$

$$(A \otimes B)X = \underbrace{\begin{pmatrix} a_{11} & & 0 & & a_{1c_A} & & 0 \\ & \ddots & & & & \ddots & \\ 0 & & a_{11} & & 0 & & a_{1c_A} \\ & & \vdots & & & & \vdots \\ a_{r_A 1} & & 0 & & a_{r_A c_A} & & 0 \\ & & \vdots & & & & \vdots \\ 0 & & a_{r_A 1} & & 0 & & a_{r_A c_A} \end{pmatrix}}_{\text{Repeating of diagonal matrices with the coefficients of } A} \underbrace{\begin{pmatrix} \overbrace{B \quad 0}^{c_A \text{ repetitions of } B} \\ \vdots \\ \underbrace{0 \quad B} \end{pmatrix}}_{\text{Block diagonal matrix}} \begin{pmatrix} x_1 \\ \vdots \\ \vdots \\ x_{c_A c_B} \end{pmatrix}$$

$\overleftarrow{\quad} r_B \quad \overrightarrow{\quad}$

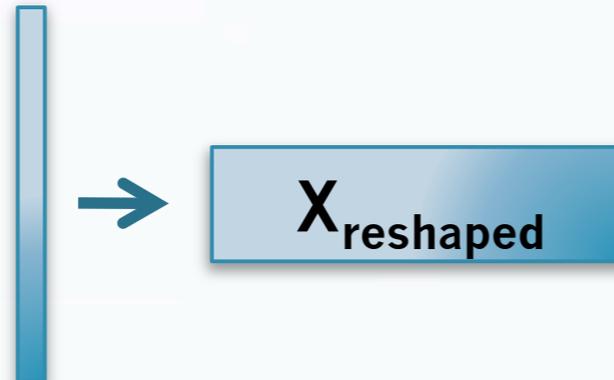
➤ Selection on the multiplied parts of X

2. KP multiplication algorithm

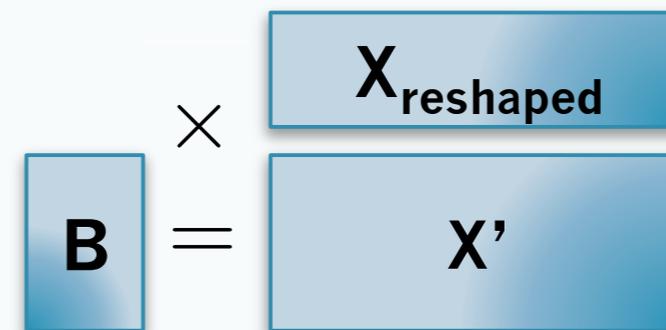
- **Successive multiplications** can be done with the underlying operators A and B

Beginning of the for loop

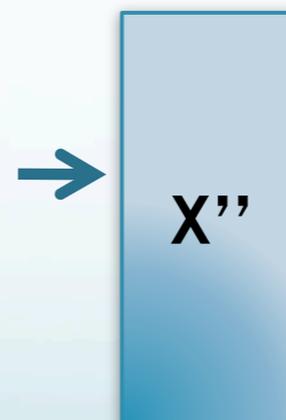
- reshaping of X



- multiplication by the operator



- Transposition of X'



End of the for loop

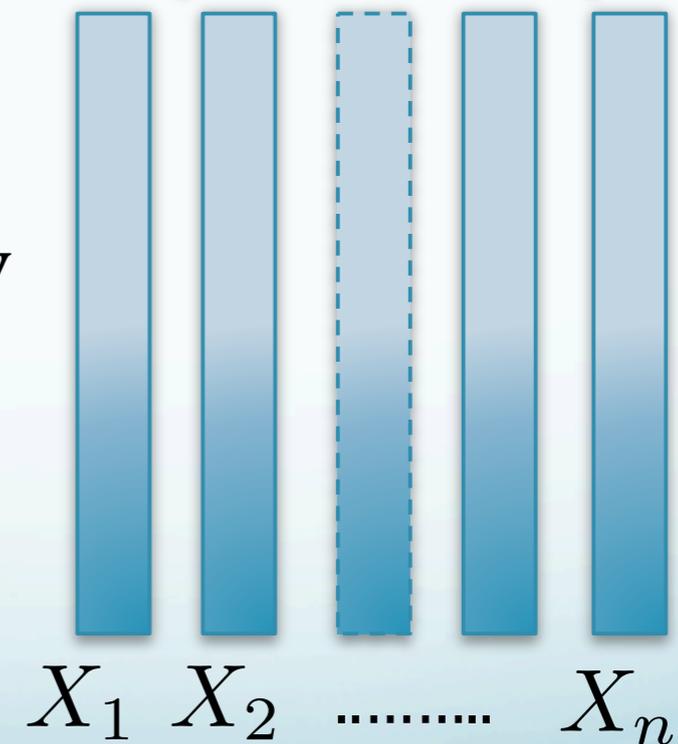
+ final collapsing

2. KP multiplication algorithm

- **If X is a matrix:** SPOT performs a *for* loop on the columns of X

$$(A \otimes B)X_1 \quad (A \otimes B)X_2 \quad \dots \quad (A \otimes B)X_n$$

- KP performed on each column, successively
- ... but, this engenders a **waste of time**



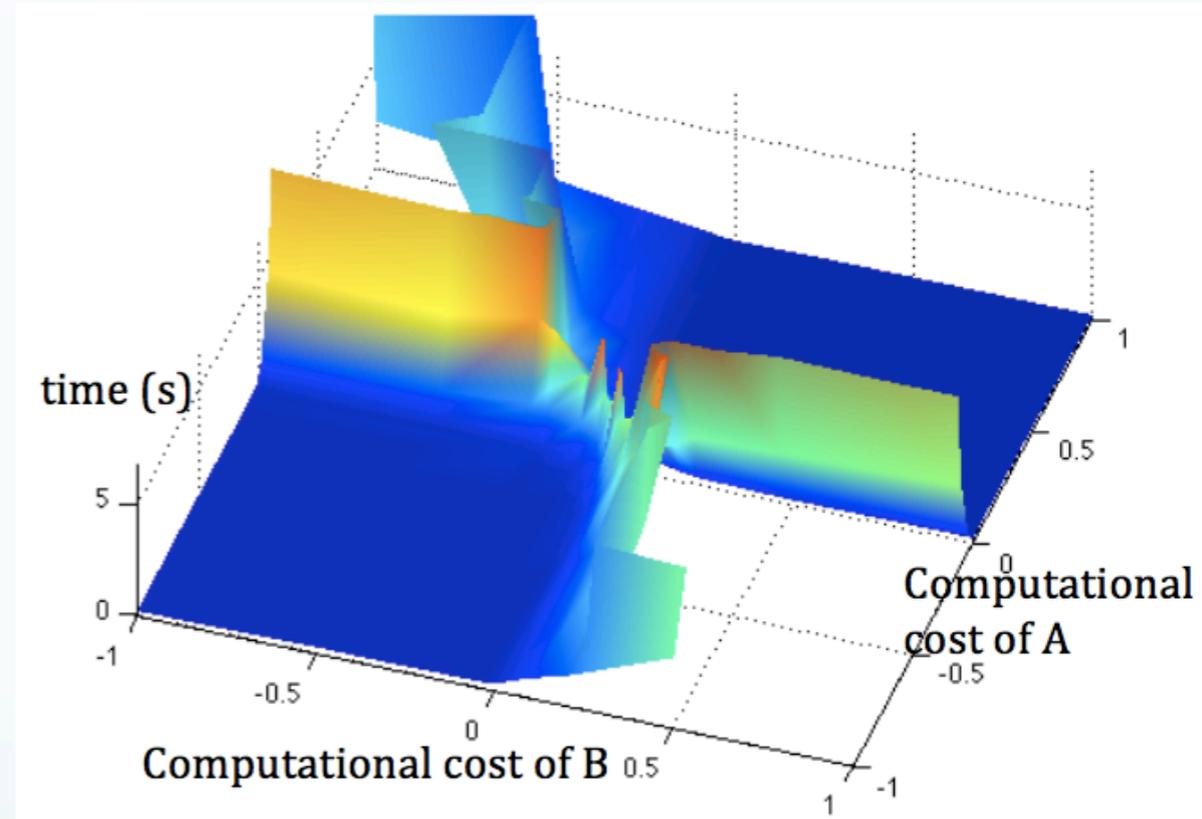
3. Initial result

➤ Computation of $(A \otimes B)X$, with X , identity matrix

- **2 high computation time zones appear**

➔ **Mainly due to non optimal computational costs**

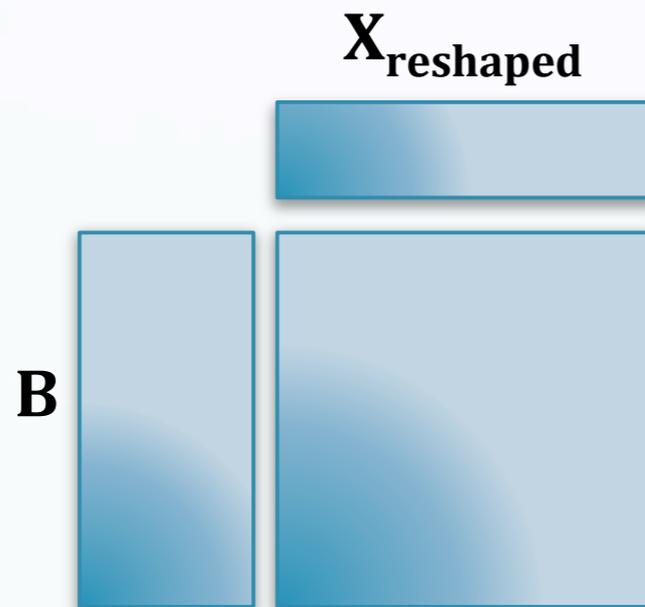
$$C = \frac{r - c}{rc}$$



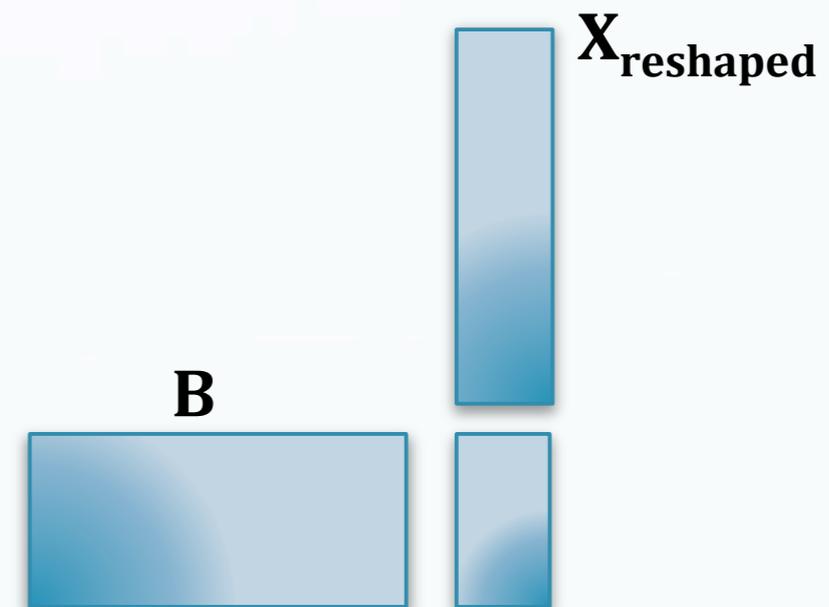
Computation time to calculate $(A \otimes B)$
 (A and B with a fixed size of 2,500 elements each)

3.Initial result

➤ Explanation



High computational cost



Low computational cost

Low computational cost \rightarrow smaller intermediate matrix

➤ Challenge

- **“to produce the smallest intermediate matrices”**

II. Kronecker product optimization

1. Permutations

- $(A \otimes B)X = (A \otimes I_{r_B})(I_{c_A} \otimes B)X = (I_{r_A} \otimes B)(A \otimes I_{c_B})X$
- More generally: **n! possibilities**

$$\begin{aligned}
 A_1 \otimes A_2 \otimes A_3 &= (A_1 \otimes I_{r_2 r_3})(I_{c_1} \otimes A_2 \otimes I_{r_3})(I_{c_1 c_2} \otimes A_3) \\
 &= (A_1 \otimes I_{r_2 r_3})(I_{c_1 r_2} \otimes A_3)(I_{c_1} \otimes A_2 \otimes I_{c_3}) \\
 &= (I_{r_1} \otimes A_2 \otimes I_{r_3})(A_1 \otimes I_{c_2 r_3})(I_{c_1 c_2} \otimes A_3) \\
 &= (I_{r_1} \otimes A_2 \otimes I_{r_3})(I_{r_1 c_2} \otimes A_3)(A_1 \otimes I_{c_2 c_3}) \\
 &= (I_{r_1 r_2} \otimes A_3)(A_1 \otimes I_{r_2 c_3})(I_{c_1} \otimes A_2 \otimes I_{c_3}) \\
 &= (I_{r_1 r_2} \otimes A_3)(I_{r_1} \otimes A_2 \otimes I_{c_3})(A_1 \otimes I_{c_2 c_3})
 \end{aligned}$$

1. Permutations

➤ Which is the best permutation ?

“The one describing operators with **increasing computational costs**”

→ Addition of a method calculating the best permutation

- Performs a quick sort
- Initializes a *best_permutation* property

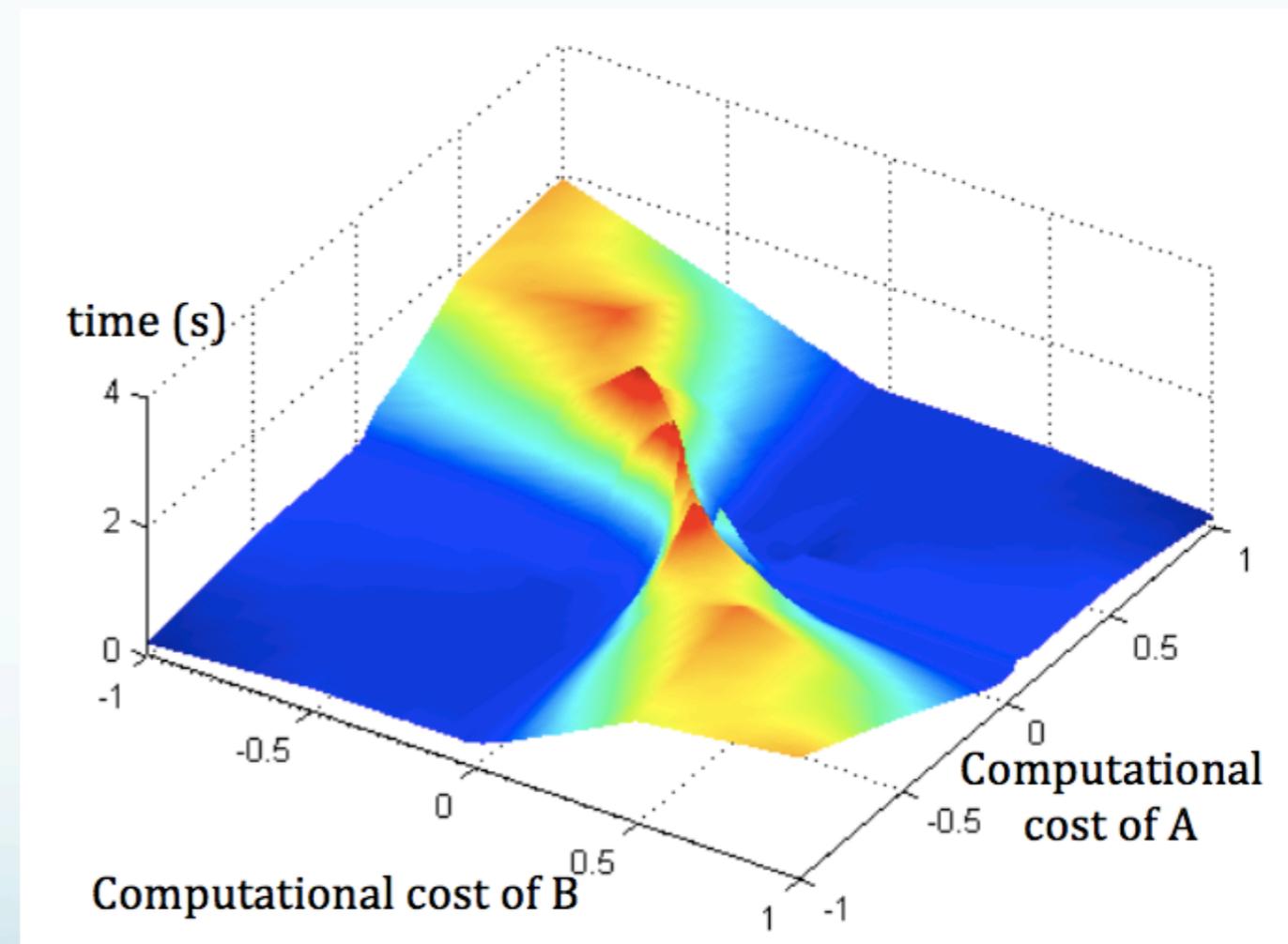
→ Multiplication according to *best_permutation* property

1. Permutations

➤ Result

Both high computation time zones have been considerably reduced

- **Average execution time**
 - divided by 12 in high computation time zones
 - about the same elsewhere



2. Replacement of the *for* loop

- Replaced by an **implicit implementation**
- *for* loop \Leftrightarrow **KP with an identity matrix**

$$I_n \otimes (\otimes_i A_i)$$

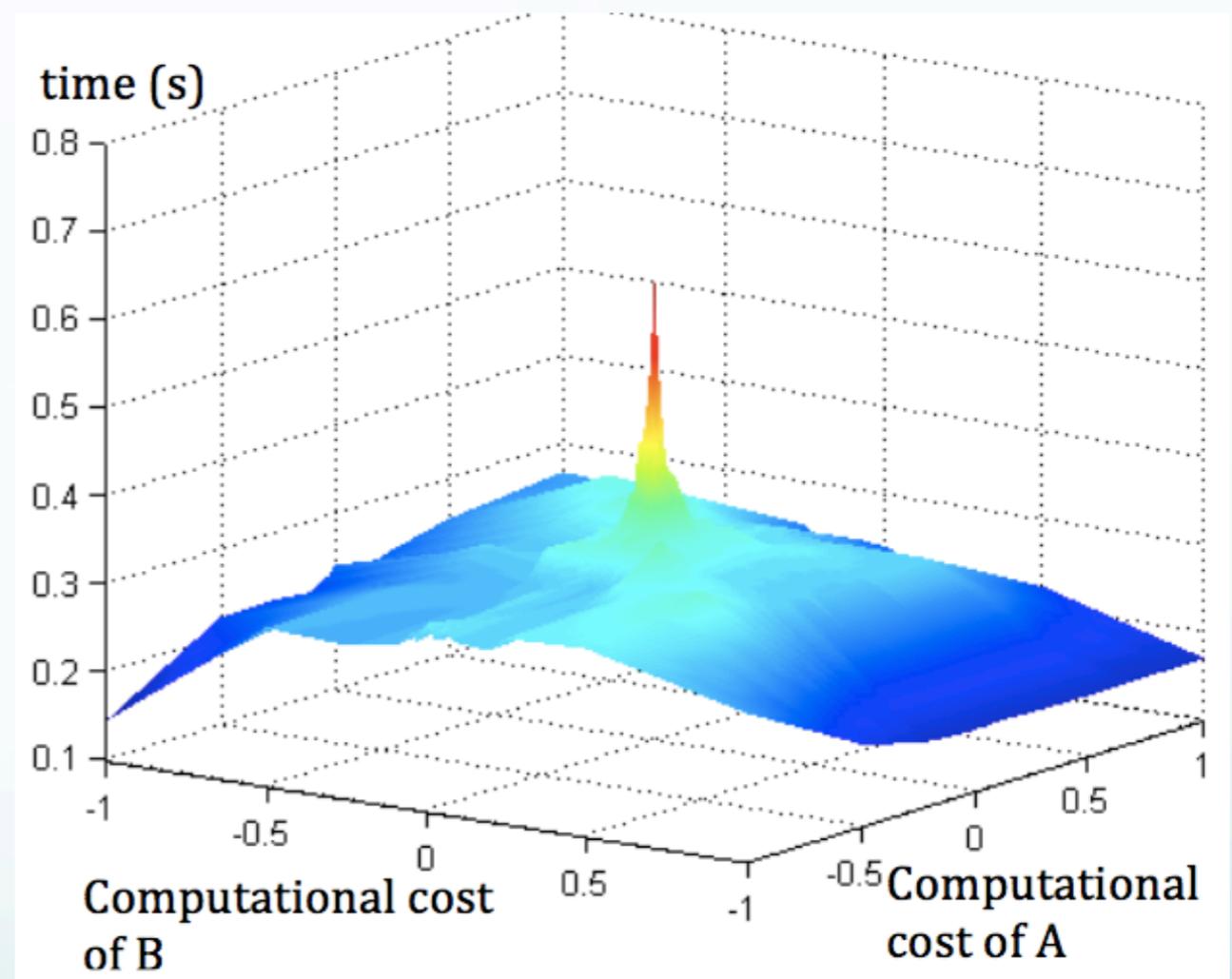
where, ***n*** : number of columns of the right-hand multiplied operator, *X*

- Equivalent to **make the size** of the first identity matrix in KP decomposition $\prod_i (I_{?} \otimes A_i \otimes I_{?})$ being ***n* times bigger**.

2. Replacement of the *for* loop

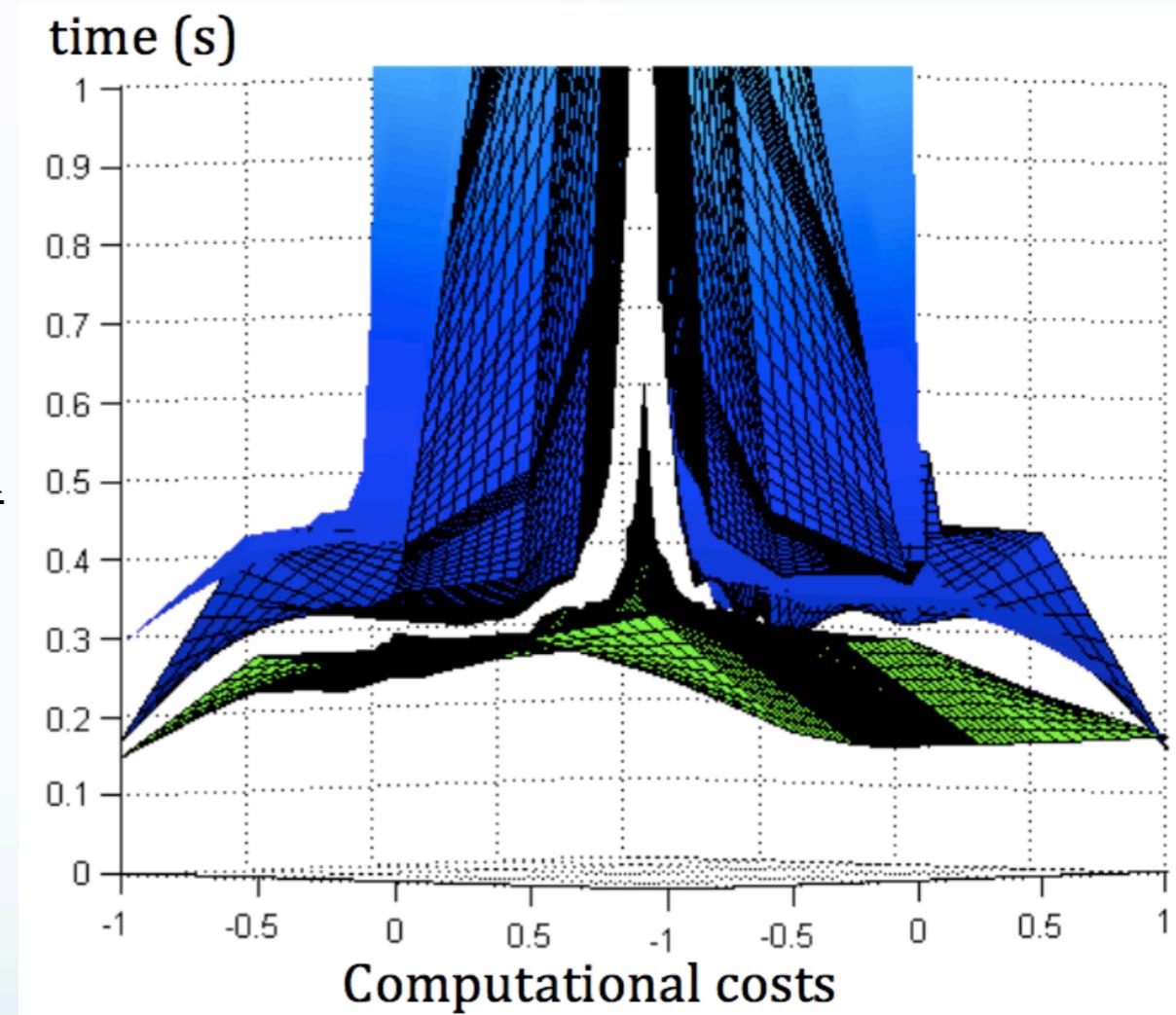
➤ Result

- **Calculation time improvement of 1,170%** compared to the initial KP multiplication function.
- **Highest computation time in the center** since it corresponds to larger intermediate matrices in calculations.



3. Final comparison

- Calculation time behavior
 - In “normal” zones
 - **divided by 2**
 - In previously high computation time zones
 - **reduced up to 50 times**
 - In the central peak region
 - **divided by 4**



Conclusion

- **Optimization of the serial Kronecker product multiplication**
 - Average gain of time of **1,170 %**
- **Improvements carried out in 2 steps**
 - Permutations → smaller intermediate matrices
 - Replacement of the *for* loop

Conclusion

- **Main advantage** of SPOT's Kronecker product
 - **KP is not entirely calculated** (consideration of its effects)
- **Next step**
 - **parallel Kronecker product** (in pSPOT)
- **Expectations**
 - **Improvements in long time processes**